



# Dynamic Pricing with Procedural and Substantive Fairness

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# Outline

- Motivation
  - Two fairness concerns
  - Doubly-fair and profitable policies
- Problem Setup
  - Online pricing with two fairness concerns.
- Algorithm
  - A policy-elimination scheme
  - Regret and Unfairness bounds

References:

Xu, Jianyu, Dan Qiao, and Yu-Xiang Wang. "Doubly Fair Dynamic Pricing." *AISTATS* 2023.

# Two Unfairnesses while Booking/Boarding Flight

- While *booking* a flight (on a 3<sup>rd</sup>-party website)...
  - Your colleague  $C$  gets a cheaper offer
  - **Proposed** prices are not equal
  - A **procedural unfairness**
- While *boarding* a flight ...
  - Your neighbor  $N$  paid at a cheaper price
  - **Accepted** prices are not equal
  - A **substantive unfairness**



# Fixed-Price Policy: A straightforward solution

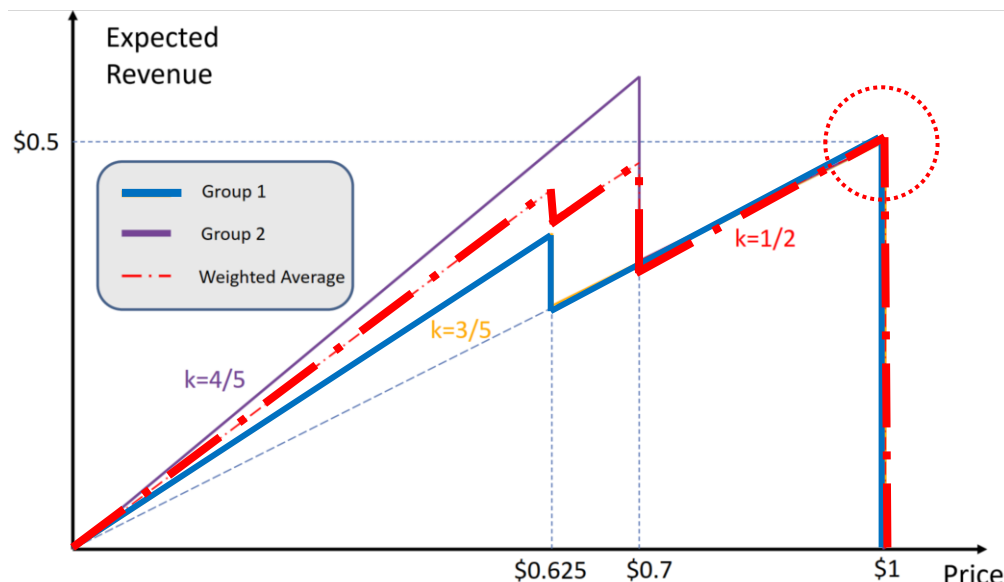
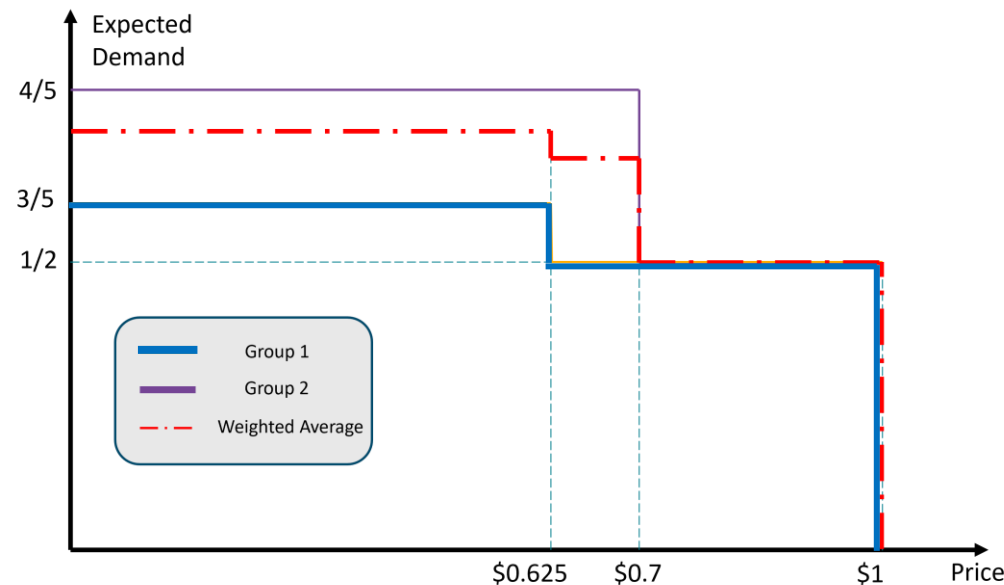
- Two fairness concerns:
  - Procedural unfairness:  $U(p_A, p_B) := |\mathbb{E}[p_A] - \mathbb{E}[p_B]|$
  - Substantive unfairness:  $S(p_A, p_B) := |\mathbb{E}[p_A | A \text{ accept } p_A] - \mathbb{E}[p_B | B \text{ accept } p_B]|$
- $p_A = p_B$  eliminate both unfairnesses.
  - Optimal price:  $p^* = \arg \max_p p \cdot (D_A(p) * Q(A) + D_B(p) * Q(B))$
  - $Q(A)$  and  $Q(B)$  are the portion of Group A and B.
- But can we do better?
  - ... if  $p_A$  and  $p_B$  are generally random.
  - Note: same distributions of  $p_A$  and  $p_B$  do not work for substantive fairness.

# Example: Randomized Prices are More Profitable

- Consider the following example:

Acceptance Rate	\$0.625	\$0.7	\$1
$G_1$ (30%)	$3/5$	$1/2$	$1/2$
$G_2$ (70%)	$4/5$	$4/5$	$1/2$

- Let  $p_1 = p_2$  to meet the fairnesses
  - $p^* = \arg \max_p D_1(p) * 0.3 + D_2(p) * 0.7 = \$1$
- If  $p_1 = p_2 = \$1$ , then
  - $D_1(p_1) = D_2(p_2) = 0.5$
  - Revenue = \$0.5
- However, if  $p_1 \sim \mathbb{P}_1, p_2 \sim \mathbb{P}_2$  are **random**, then ...



# Example: Randomized Prices are More Profitable

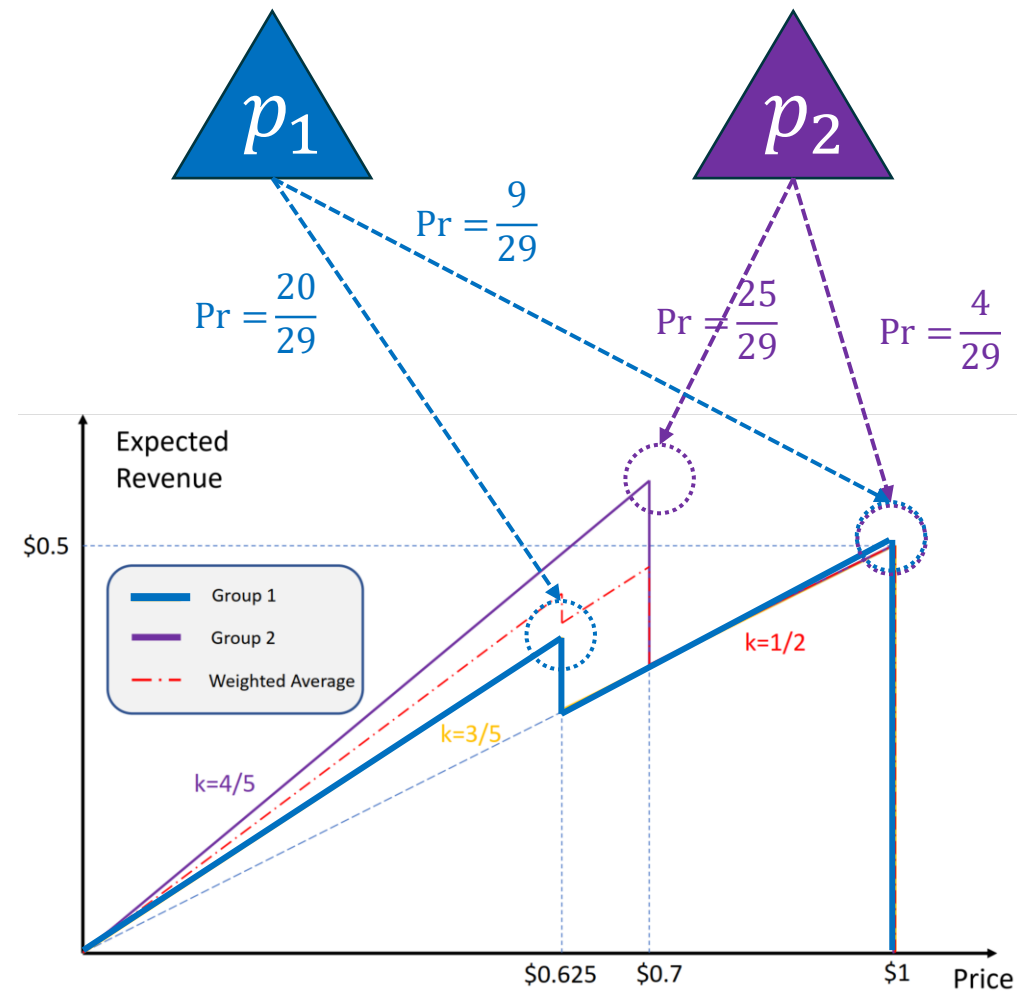
- Let  $p_1 = \begin{cases} \$ 0.625 & (\text{Pr} = \frac{20}{29}) \\ \$ 1 & (\text{Pr} = \frac{9}{29}) \end{cases}$  and  $p_2 = \begin{cases} \$ 0.7 & (\text{Pr} = \frac{25}{29}) \\ \$ 1 & (\text{Pr} = \frac{4}{29}) \end{cases}$

- Procedural fairness holds as  $\mathbb{E}[p_1] = \mathbb{E}[p_2] = \frac{43}{58}$ .

- Substantive fairness holds as  $\mathbb{E}[p_1 | \text{buy}] = \mathbb{E}[p_2 | \text{buy}] = \frac{8}{11}$

- And **profit increases !!**

- $\mathbb{E}[p_1 \cdot 1[p_1 \text{ accepted}]] \cdot 0.3 + \mathbb{E}[p_2 \cdot 1[p_2 \text{ accepted}]] \cdot 0.7$
- $= \$ \frac{74}{145} \approx \$0.5103 > \$ 0.5.$



# Seek for Optimal Price Distribution

- Find optimal price  $\rightarrow$  optimal *distribution* of prices.
- Question 1: What is the best fair distribution?

$$\begin{aligned} \max_{\pi_1, \pi_2} R(\pi_1, \pi_2) &:= \mathbb{E}_{p_1 \sim \pi_1} [p_1 \cdot D_1(p_1)] \cdot Q_1 + \mathbb{E}_{p_2 \sim \pi_2} [p_2 \cdot D_2(p_2)] \cdot Q_2 \\ \text{s. t. } U(\pi_1, \pi_2) &= 0, \quad S(\pi_1, \pi_2) = 0 \end{aligned}$$

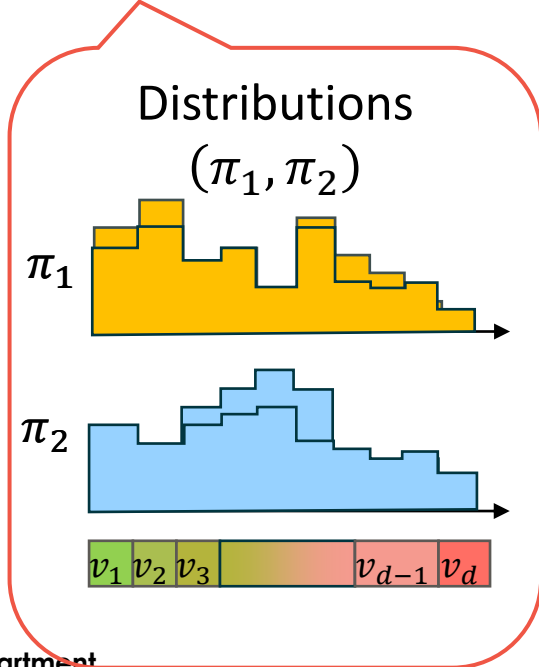
Here  $D_i(p)$  are the demand function of each group  $i$ .

- Question 2: How to **learn** them over time?
  - Unknown  $D_1(p)$  and  $D_2(p) \rightarrow$  Unknown  $R(\pi_1, \pi_2)$  and  $S(\pi_1, \pi_2)$
  - But we can learn from customers' buy/not buy decisions.
  - An online dynamic pricing problem.

# Online Pricing: Problem Setup

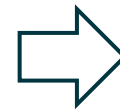
At time  $t + 1$ :

Seller



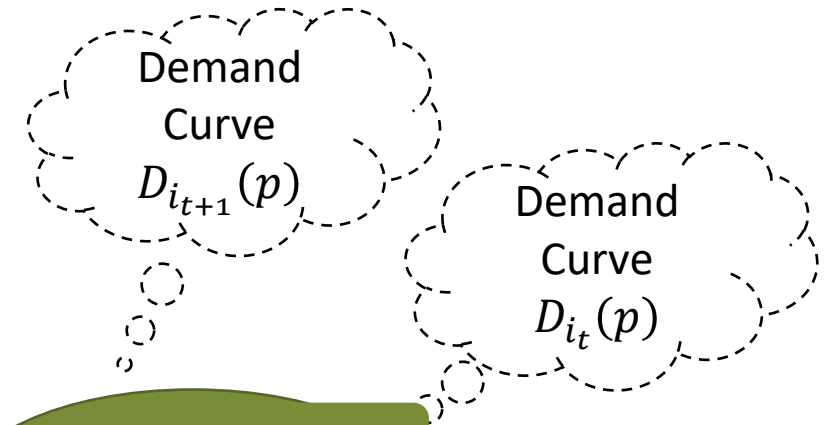
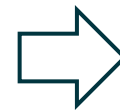
$\sim$

$p_t$



$\sim$

$p_{t+1}$



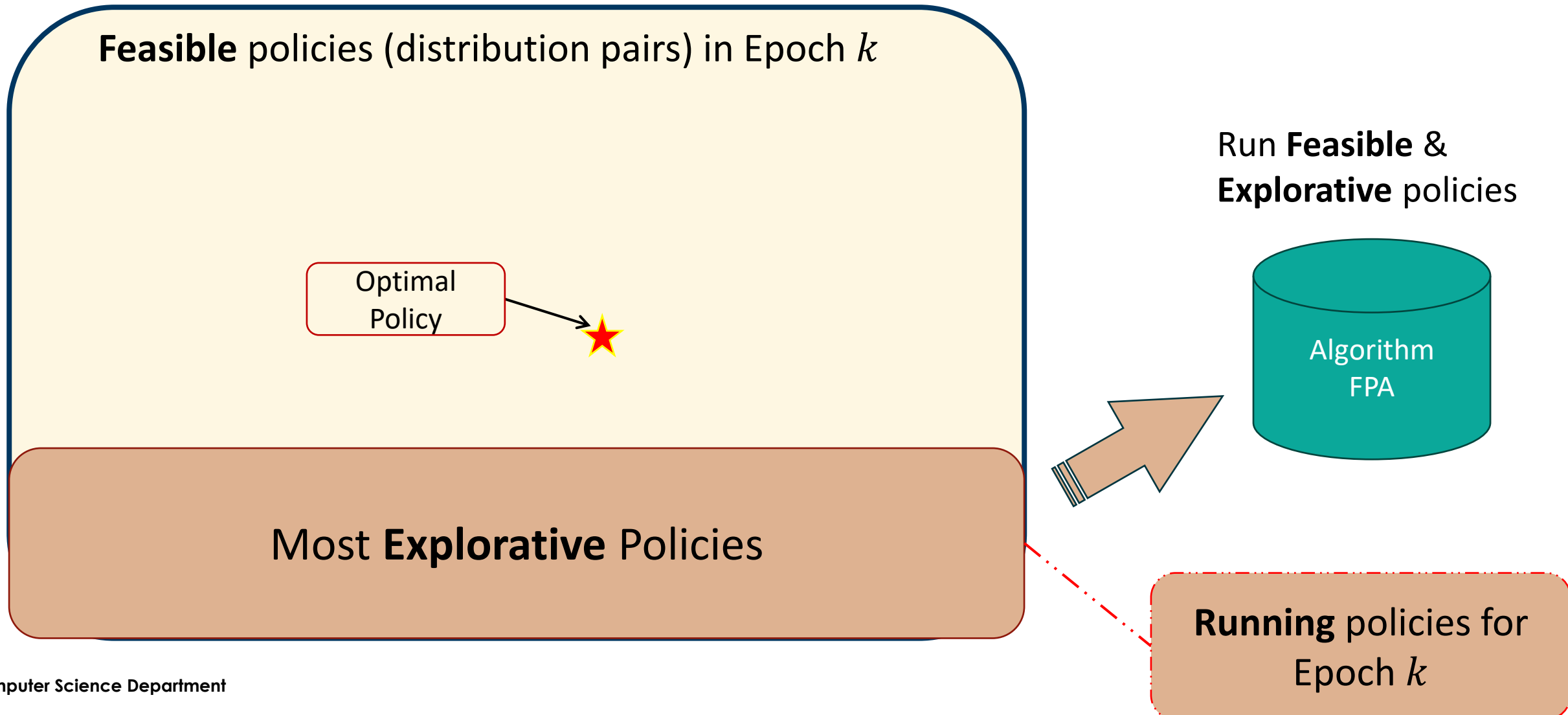
Customer  $t+1$   
From Group  $i_{t+1}$

Decision  
 $1_t \sim \text{Ber}(D_{i_t}(p_t))$

Decision  
 $1_{t+1} \sim \text{Ber}(D_{i_{t+1}}(p_{t+1}))$



# Algorithm: Epoch-based Policy Elimination



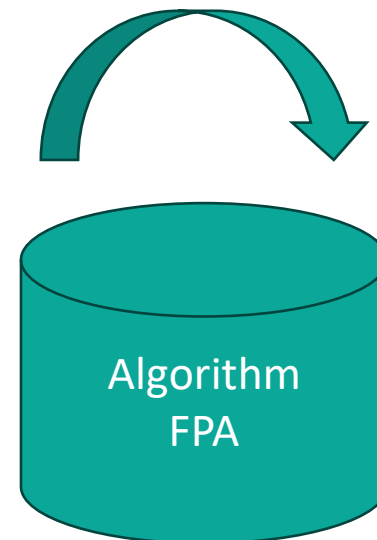
# Algorithm: Epoch-based Policy Elimination

Feasible policies (distribution pairs) in Epoch  $k$

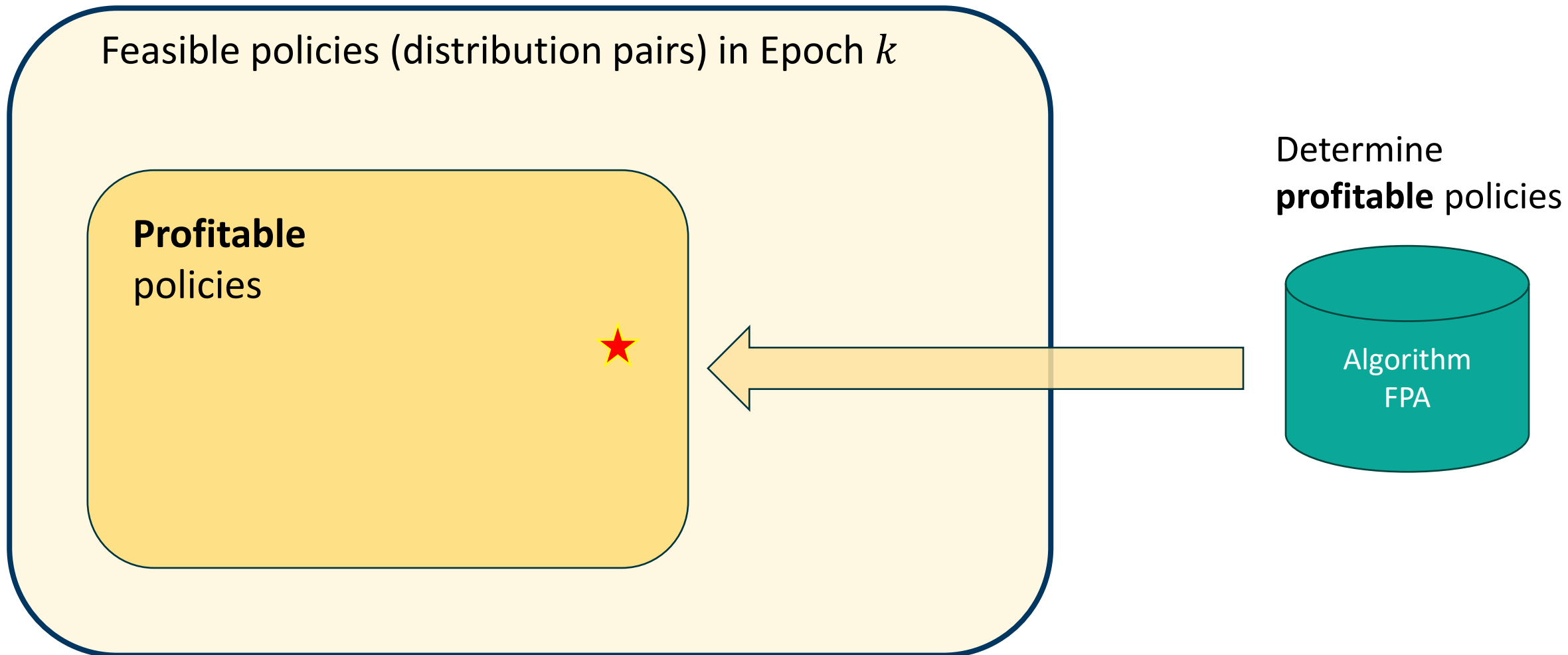


Most **Explorative** Policies

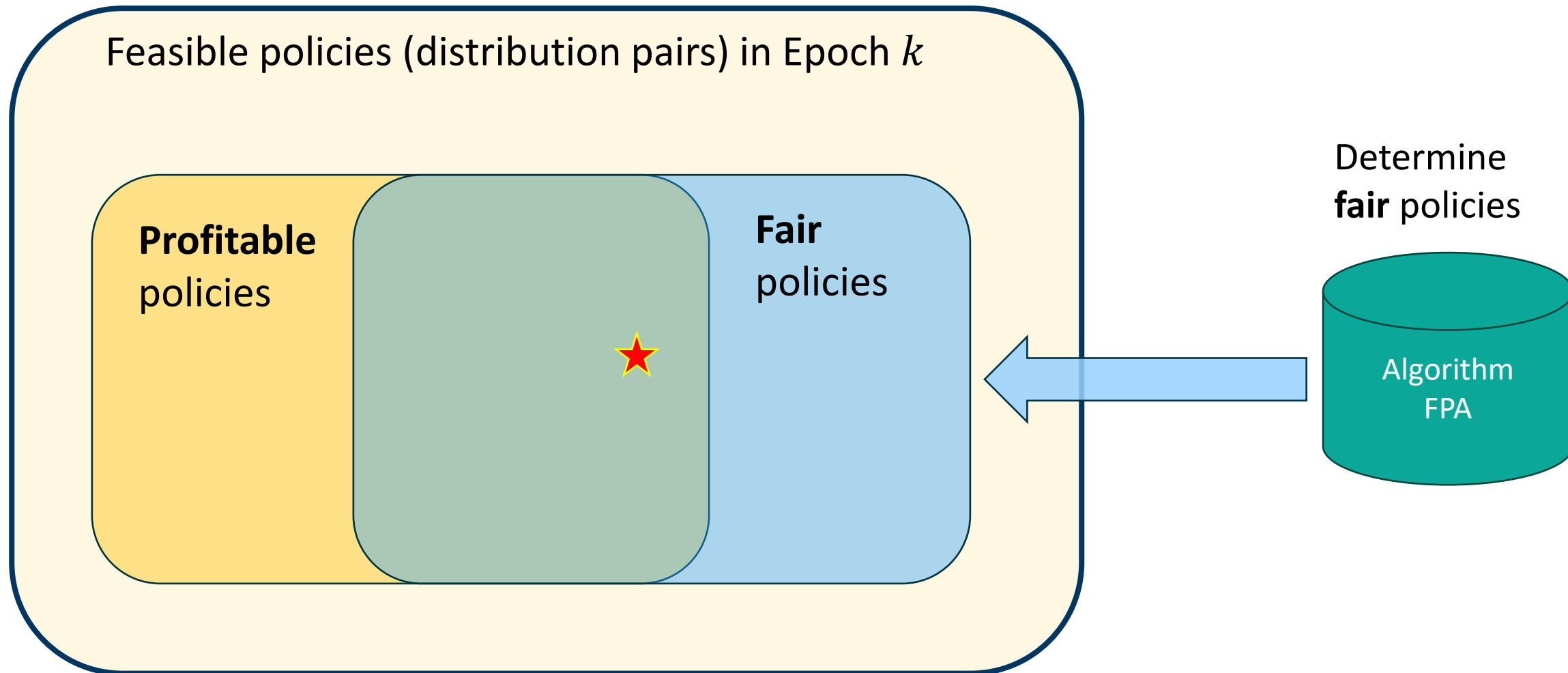
Gather feedback and solve constraint optimization problems



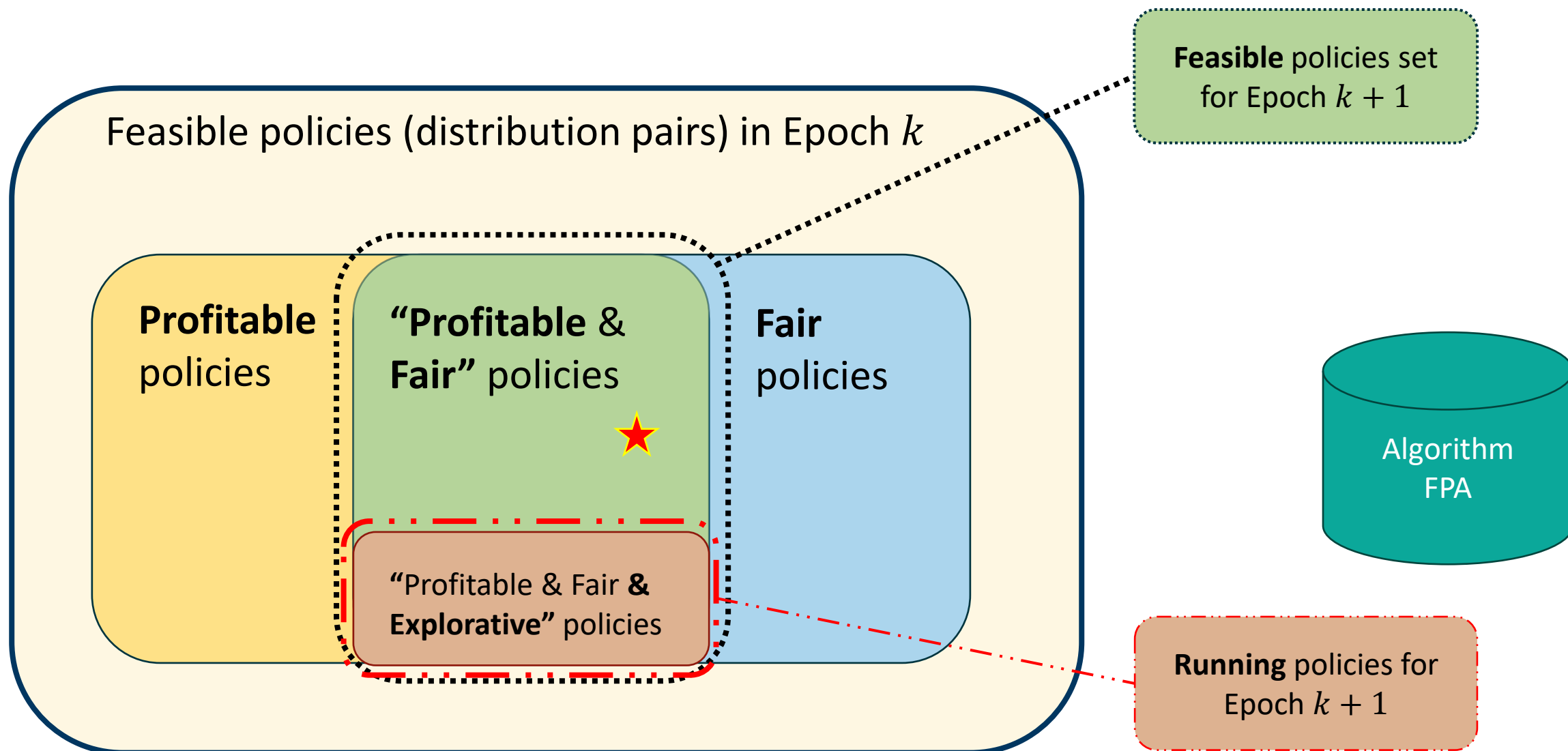
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# Algorithm: Epoch-based Policy Elimination



# Optimal Regret, Optimal Unfairness, and Optimal Regret-Unfairness Tradeoffs

- Our FPA algorithm guarantees ...
  - $\tilde{O}(\sqrt{T})$  regret
  - **0** procedural unfairness
  - $\tilde{O}(\sqrt{T})$  substantive unfairness
- To show the optimality, we also prove lower bounds of ...
  - $\Omega(\sqrt{T})$  regret
  - Necessarily unfair:  **$O(\sqrt{T})$  regret  $\Rightarrow \Omega(\sqrt{T})$  substantive unfairness.**

# Conclusions and Potential Extensions

- Two fairness concerns:
  - Procedural and Substantive fairness
- Randomized prices might be doubly-fair and more profitable
- To solve the online fair pricing problem, we
  - Propose an FPA algorithm
  - Prove its regret and unfairness guarantees
  - Show the optimality of regret & unfairness & regret-unfairness tradeoffs.

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